Python based implementation of a chaos based cryptosystem using external key

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Outline of the presentation

- *My* introduction to Python
- Introduction to Cryptography and Chaos
- The paper's core
- Results
- Conclusions
- References

My Introduction to Python

- Internet
- Need for a commonly used scripting language
- Open source/free
- Lots of documentation
- Well connected
- Mature

Introduction to Cryptography

- Ever increasing trend towards virtualization
- Various valuable transactables like information, education & money
- Continuous research towards newer cryptographic implementations
- Rapid progress in Hardware
- Rapid progress in Software

Introduction to Chaos

- Pseudo-random number generation $X_{n+1} = f(X_n)$
- For e.g., the logistic map: $X_{n+1} = pX_n(1 X_n)$, p = 3.99
- High sensitivity to initial conditions & system parameters
- Wide variety of chaotic maps
- Multidimensional chaotic maps
- Discrete/continuous maps
- Applications to cryptography [1-3]

Introduction to the Chaos based Cryptosystem - overview

• Two skew-tent chaotic maps [2]:

 $t(x) = x/p, when 0 \leq x < p$ $t(x) = (1-x)/(1-p), when p \leq x \leq 1$

- The two maps differ in terms of the parameter *p* and the number of iterations for each new random number
- A 128-bit key is used to initialize the chaotic maps and to specify other operational parameters

• Plaintext is broken into 8-byte blocks as:

 $P = P_1 P_2 P_3 \dots P_m$

The last block may be smaller

• The resulting ciphertext is similarly obtained in the form of 8-byte blocks:

 $C = C_1 C_2 C_3 \dots C_m$

- The two skew-tent maps have p = 0.79 & 0.39 resp.
- The first map has a variable no. of iterations T₀ and the second map has only one iteration per step

- The 128-bit key is broken into four 32-bit parts: $K=A_1A_2B_1B_2$
- The initial value of T_0 is an 8-bit integer, obtained as, $T_0 = (lowest \, 4 \, bitsof A_2) (highest four bitsof B_1)$
- The initial values X_0 and Y_0 are: $X_0 = 0.((A_1A_2)XOR(B_1B_2))$
 - $Y_0 = 0.(A_1 B_2)$

- Using these initial values (X₀, Y₀ and T₀), the two tent-maps are run to the updated X_b and Y_b
- The next-to-decimal 64 bits of X_{b} and Y_{b} are used as random integers X_{bz} and Y_{bz} in the further operations
- Block-wise encryption and decryption are defined as: $C_b = s(P_b) XORC_{b-1} XORX_{bz} XORY_{bz}$

 $P_{b} = s^{-1} \left(C_{b} XORC_{b-1} XORX_{bz} XORY_{bz} \right)$

- Here, $C_0 = A_2 B_1$
- The *s*(..) operator consists of the following 2 steps:
- 1. Each byte of the block P_{b} is rotated to a new position in the block P_{b}' the rotation number being the byte-wise sum of the block modulo the block-size
- 2. The bits of of the resulting block P_b' are swapped left-right – the swap location being the 1-bit population count of $A_1 B_2$

- The tent-maps' iterations are re-initialized as follows: $X_{bz} = C_{b-1} XORX_{(b-1)z} XORB_2 A_1$ $Y_{bz} = C_{b-1} XORY_{(b-1)z}$ $T_b = z (P_{b-1}) XORT_{(b-1)}$
- The *z*(..) operator being the byte-wise XOR operator
- The above steps are repeated until the complete plaintext/ciphertext is exhausted

Application of gmpy

- GMP GNU Multi-precision library
- gmpy the *friendly* Python interface to the gmp

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• Example code:
import random, gmpy
mykey1 = gmpy.mpz(random.randrange(0, 2**128 - 1))
myB = mykey1.lowbits(64)
myA = mykey1 - myB
myA = myA.__rshift__(64)
myT0 = myA.lowbits(4)
myT0 = myT0.__lshift__(4)
myT0 = myT0 + myB.__rshift__(60)
```

The commands

The encryption:

./myxiang08v2.py e inaroundcampus.pdf goodpw1
The decryption:
./myxiang08v2.py d einaroundcampus.pdf goodpw1

The histogram:

./myhistogram.py inaroundcampus.pdf

./myhistogram.py einaroundcampus.pdf

Results



Results – contd.



Conclusions

- Python, as often stated, was found to be extremely simple and user-friendly
- gmpy provided a friendly interface to an otherwise complicated C-library
- The actual Python code was extremely short in comparison to the C-code
- The code was easily understandable and reusable
- Several straightforward extensions are possible biometrics, audio/video encryption etc.

References

[1] Muhammad Usama , Muhammad Khurram Khan , Khaled Alghathbar and Changhoon Lee, *Chaos-based secure satellite imagery cryptosystem*, Computers and Mathematics with Applications 60 (2010) 326–337.

[2] Tao Xiang, Kwok-wo Wong and Xiaofeng Liao, An improved chaotic cryptosystem with external key, Communications in Nonlinear Science and Numerical Simulation 13 (2008) 1879-87.

[3] Pareek NK, Patidar V, Sud KK. *Cryptography using multiple one-dimensional chaotic maps*, Communications in Nonlinear Sciences and Numerical Simulations, 10 (2005), 715-23.

THANKS! Have a nice day!